## BIKE - Bit-Flipping Key Encapsulation

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- McEliece-like KEM with QC-MDPC Codes
- Well-Understood & Reliable Security
  - Theoretical Security: Reduction based on well-known coding-theory problems
  - Practical Security: ISD-based attacks [Pra62] whose work-factor\* barely changed in ~50 years of research
- Performance
  - **Practical performance** for all KeyGen, Encaps, Decaps steps regarding both computational complexity and bandwidth

## Tip: Ideal usage as Ephemeral Key Exchange (e.g. SSL/TLS)

# NIST Report 8240 on 1<sup>st</sup> Round BIKE



"BIKE targets IND-CPA security and makes no attempt to make it difficult for an attacker to mount a chosen ciphertext attack if keys are reused. This design decision was made by the submitters, based on the difficulty of designing a bit-flipping decoder with a low enough decoding failure rate to allow an efficient IND-CCA2-secure construction."

NIST IR Report 8240, page 11, Section 3.12

## New Backflip Decoder



- Context
  - Round 1 decoder: efficient but fails with non-negligible probability (10<sup>-7</sup>)
  - To enable IND-CCA variants, negligible decoding probability was needed
- Backflip Rationale
  - Similar to Bit-Flipping
  - <u>Difference</u>: each bit flip keeps a time-to-live counter. After a given number of iterations, the bit flip reaches a time-to-death point and is flipped back
  - <u>Result</u>: Based on an extrapolation argument, it is possible to show that a certain parameter set attains an arbitrarily low failure rate using Backflip

## New BIKE CCA Variants



- Core Ingredients
  - Backflip Decoder
  - [HHK17] like conversions (with bounds from [JZC18, JZM19])
- CPA→CCA conversion preserved the strong points of each variant

	Strong Points among CPA Variant	Strong Points among CCA Variant
BIKE-1	Fastest KG+Encaps+Decaps among CPA variants	Fastest KG+Encaps+Decaps among CCA variants
BIKE-2	Smallest ciphertext among CPA variants	Smallest ciphertext among CCA variants
BIKE-3	Security reduction to single problem	Security reduction to single problem

• CCA Variants enable static keys. Current focus remains CPA Ephemeral

# NIST Report 8240 on 1<sup>st</sup> Round BIKE



## "BIKE offers key and ciphertext sizes and performance that are competitive with ring and module lattice schemes (especially at the lower security categories)."

NIST IR Report 8240, page 11, Section 3.12

BIKE CCA -- constant time implementations Nir Drucker <sup>1, 2</sup>, Shay Gueron <sup>1, 2</sup>, Dusan Kostic <sup>1, 3</sup> (1) Amazon (2) University of Haifa (3) EPFL

- 3
- New BIKE CCA implementation in constant time -- C, AVX2, AVX512
  - Constant time algorithm definition for the CCA decoder
  - Constant time implementation for the CCA BIKE flows
- Conclusions
  - It is possible to define and implement BIKE CCA in constant time
  - Performance costs are tolerable
- "Additional" code package & detailed report to be released/published soon

## BIKE CCA -- constant time implementations Nir Drucker <sup>1, 2</sup>, Shay Gueron <sup>1, 2</sup>, Dusan Kostic <sup>1, 3</sup> (1) Amazon (2) University of Haifa (3) EPFL

		CPA Cycles	CCA Cycles	CCA/CPA	CPA Bandwidth	CCA Bandwidth	CCA-CPA
BIKE 1	KeyGen	250K	270K	1.08x			
	Encaps	180K	210K	1.17x	→ 2.48KB ← 2.48KB	→ 2.88KB ← 2.88KB	<b>Δ</b> =400B
	Decaps	1.9M	8.1M	4.18x			
BIKE 2	KeyGen	4.8M	6.3M	1.30x	→ 1.24KB	→ 1.44KB	
	Encaps	140K	170K	1.21x	← 1.24KB	← 1.47KB	<b>Δ</b> =400B
	Decaps	1.9M	9.0M	4.64x			
BIKE 3	KeyGen	220K	250K	1.14x	→ 2.69КВ	→ ЗКВ	
	Encaps	220K	280K	1.27x	← 2.69KB → 1.38KB	← 3.03KB → 1.53KB	<b>Δ</b> =310B
	Decaps	2.5M	13.8M	5.58x	← 2.69KB	← 3.03KB	

### Parameter sets targeting NIST Security Level 1

In red message size for BIKE-3 with compressed g

# s2n is an AWSLabs open source library for TLS Small, fast, with simplicity as a priority Removes a lot of cruft that has built-up in libssl

BIKE Real Experiment with

Currently handles all of the S3 traffic today

### • PQ-TLS 1.2 – hybrid key exchange in s2n

s2n AWS TLS library

- Added SIKE and BIKE (reference code) into the s2n code base
- Added a hybrid key exchange cipher suites into s2n
  - TLS\_ECDHE\_BIKE\_RSA\_WITH\_AES\_256\_GCM\_SHA384
  - TLS\_ECDHE\_SIKE\_RSA\_WITH\_AES\_256\_GCM\_SHA384
- Applied the same rigor to this new code as in all of s2n
- Open Source implementation will be released soon

### Conclusion: feasible to use "Classical + BIKE" hybrid in a real networking application



# NIST Report 8240 on 1<sup>st</sup> Round BIKE



"Security strengths are based on information-set-decoding attacks, which have a long history of analysis during which the complexity of such attacks have not greatly changed.

NIST IR Report 8240, page 11, Section 3.12

"Possible areas for further analysis related to BIKE include ... investigating the effect, if any, of the quasi-cyclic code structure on security."

NIST IR Report 8240, page 12, Section 3.12

# The Effect of Quasi-Cyclic Code Structure on Security



- QC-MDPC Parameters are selected considering three ISD-related attacks
  - Key distinguishing attack: Exhibit one codeword of  $C^{\perp}$  of weight w
  - Key recovery attack: Exhibit r codewords of  $C^{\perp}$  of weight w
  - **Decoding attack**: Decode t errors in a (n, n r)-linear code.
- ISD algorithms assume a list of solution candidates of size L. Each candidate has a probability P to produce a solution. Under optimal conditions:  $WF_{ISD}(n,r,t) \approx L/P$
- [Sen11] shows that the gain when the decoding problem has  $N_s$  solutions and when  $N_i$  instances are treated simultaneously is:  $N_s/\sqrt{N_i}$

	MDPC	QC-MDPC
Key distinguishing	$\frac{1}{r} \mathrm{WF}_{\mathrm{isd}}(n, n-r, w)$	$\frac{1}{r}$ WF <sub>isd</sub> $(n, n - r, w)$
Key recovery	$\operatorname{WF}_{\operatorname{isd}}(n, n-r, w)$	$\frac{1}{r}$ WF <sub>isd</sub> $(n, n - r, w)$
Decoding	$\mathrm{WF}_\mathrm{isd}(n,r,t)$	$\frac{1}{\sqrt{r}} WF_{isd}(n, r, t)$

See [Mis13] for a detailed analysis

# Smaller Updates & Final Remarks



## • Smaller updates

- BIKE-3 variant that generates g from a seed, saving almost 50% communication
- Fixed decoding threshold computation in reference & optimization code, which now matches the spec, accelerating decoding. No changes in additional code;
- Fixed buffer overflows in reference & optimization code;
- Final remarks
  - BIKE has well-understood, reliable security & practical performance
  - BIKE is particularly appealing for low-level security (e.g. Level 1)
  - Given CPA focus, variants with fast key generation (e.g. BIKE-1, BIKE-3) are our priority
  - NIST Report 8240 already highlights benefits of BIKE and the team addressed requests





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